





Faster Multi-Object Segmentation using Parallel Quadratic Pseudo-Boolean Optimization

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Summary

In this paper we:

Introduce a parallel algorithm for solving quadratic pseudo-Boolean optimization (QPBO) problems.

Speed-up Results

Data & Setup

We apply our algorithm to two different tasks:



Nerves



- Demonstrate an 11x and 19x speed-up for large problems and ~3x for small problems.
- Demonstrate our method outperforms current state of the art algorithms by a large margin.

We focus on image segmentation, but our algorithm can solve any QPBO problem. Our implementation is available at: github.com/Skielex/shrdr

Method

Graph Construction

Given a quadratic pseudo-Boolean energy function

$$E(\mathbf{x}) = \sum_{p \in \mathcal{V}} \theta_p(x_p) + \sum_{p,q \in \mathcal{V}} \theta_{pq}(x_p, x_q)$$

a graph is constructed as follows



1. Nerves (N1, N2): segment nerves from a large 3D μ CT image.

2. Broad Bioimage Benchmark (BBB): segment cells from many small 2D images.

All tests are run on a dual socket machine with two 16 core Intel Xeon 6226R CPUs and enough RAM to keep all data loaded in memory.

Broad Bioimage Benchmark





	N1	N2
Nodes	364 M	818 M
Edges	2,124 M	4,864 M
Memory footprint		
K-QPBO	134.2 GB	306.8 GB
M-QPBO	60.1 GB	182.7 GB

Min-cut/max-flow of this graph gives the minimizer of the energy function.

Parallel QPBO Algorithm

We combine the original QPBO algorithm with the block-based parallelization approach by Liu and Sun. The algorithm has two phases:

- A. Split the problem into disjoint blocks and solve each block in parallel.
- B. Merge pairs of blocks and re-solve. Repeat until all blocks are merged.

During both phases, graph arcs for non-submodular terms are only added when needed. This is key to our algorithm's performance.



Results on Broad Bioimage Benchmark (Small Problems)



Performance Comparison

We compare our algorithm with:

- Liu-Sun: parallel min-cut/maxflow solver.
- M-QPBO 1.32 1.72±0.28 60.1 GB **EIBFS:** state of the art serial min-10.17 2.96±0.89 P-QPBO (32) 70.0 GB cut/max-flow solver.
- 70.0 GB 4.42 0.78±0.13 Liu-Sun (32) **P-EIBFS:** parallelized version of \bullet EIBFS 175.1 GB 0.92 0.33±0.08 EIBFS.
- P-EIBFS (32) 175.8 GB 0.59 0.68±0.14 **EIBFS-QPBO:** estimate of QPBO EIBFS-QPBO 175.1 GB 1.83 0.66±0.08 with EIBFS.

*Tested with up to 32 threads.

N1 Mem.

Best speed-up

N1

BBB